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(54) **SYSTEM AND METHOD FOR ROUTING  
DATA TO DEVICES WITHIN AN  
INFORMATION HANDLING SYSTEM**

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G06F 15/16; G06T 1/20; H04W 40/00  
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(56)

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<b>H04W 40/00</b>	(2009.01)
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<b>H04L 12/801</b>	(2013.01)
<b>H04L 12/773</b>	(2013.01)

(52) **U.S. Cl.**

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(2013.01); **H04L 47/13** (2013.01); **H04W**  
**40/00** (2013.01); **H04L 45/60** (2013.01)

(58) **Field of Classification Search**

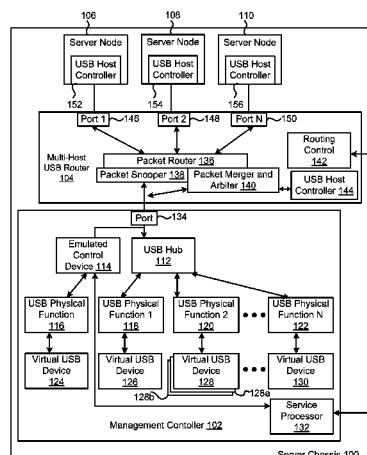
CPC ..... H04L 45/58; H04L 45/70; H04L 47/125;  
H04L 45/54; H04L 45/60; H04L 47/13;

(57)

**ABSTRACT**

An information handling system includes a management controller and a router. The management controller includes a universal serial bus hub and first and second devices. The management controller to assign the first device to a first server node, to assign the second device to a second server node, and to create a routing table associated with the assignment of the first and second devices respectively to the first and second server nodes. The router is in communication with the management controller. The router to receive the routing table from the management controller, to receive data from the first server node, and to route the data to the first device based on the routing table. The universal serial bus hub communicates with the router via a single physical port of the management controller.

**20 Claims, 5 Drawing Sheets**



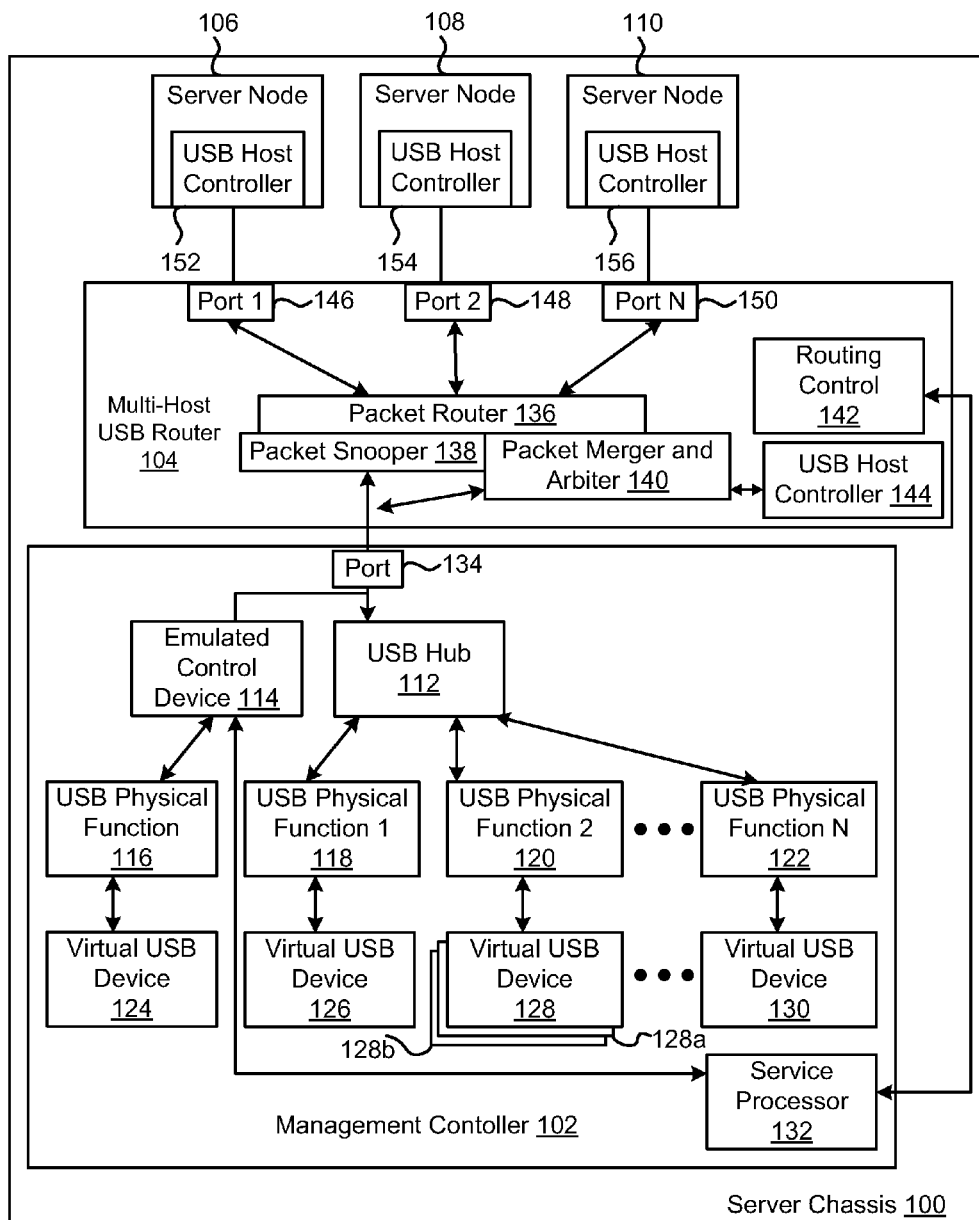
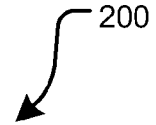


FIG. 1



Routing Table			<u>200</u>
BMC Function #	Host Port #	Device	
1	1	KVM (KB/Mouse)	202
2	2	KVM (KB/Mouse)	204
3	2	Virtual CD	206
4	3	Lifecycle Controller	208
5	4	Virtual USB Mass Storage Device	210
6	3	Virtual USB Mass Storage Device	212

**FIG. 2**

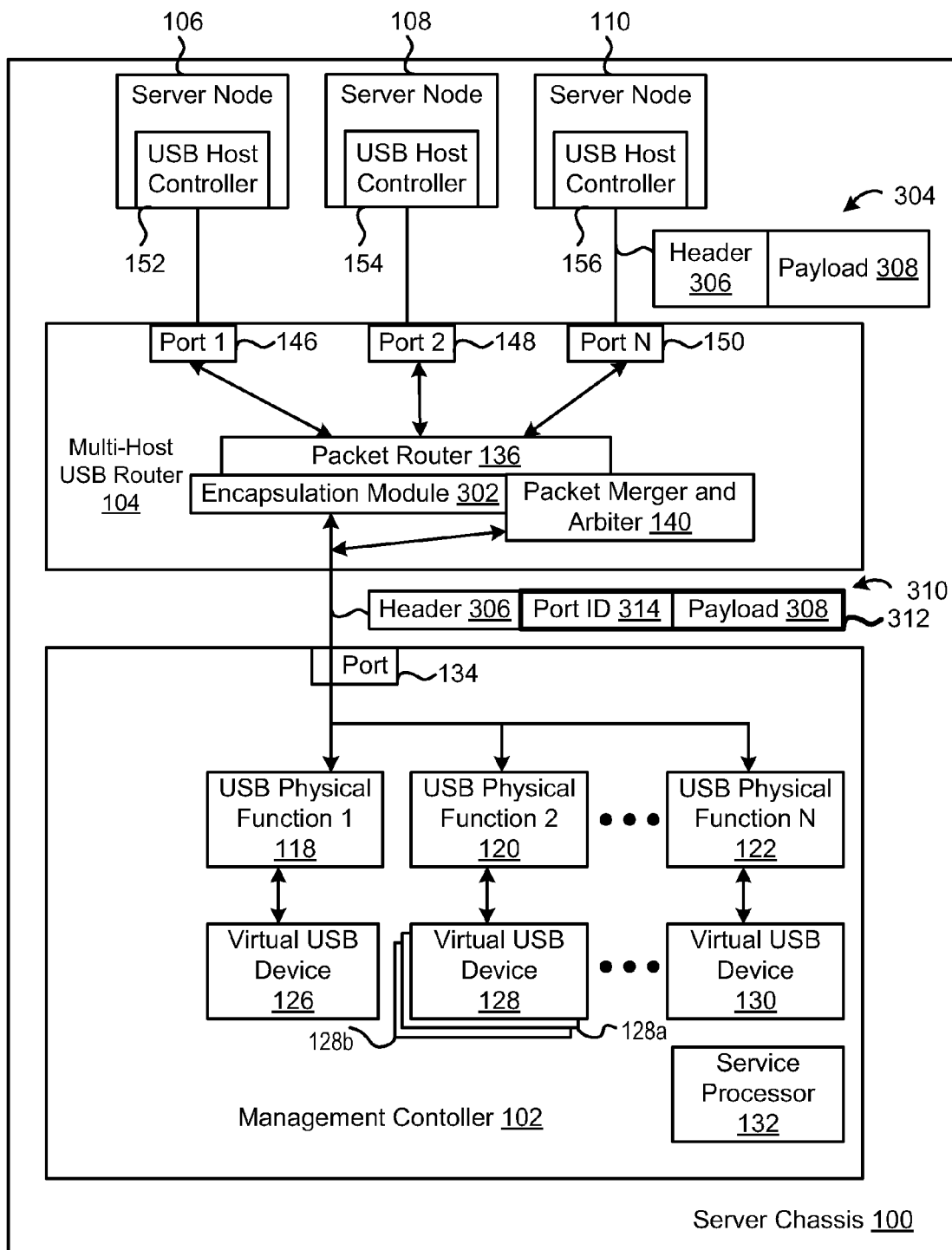
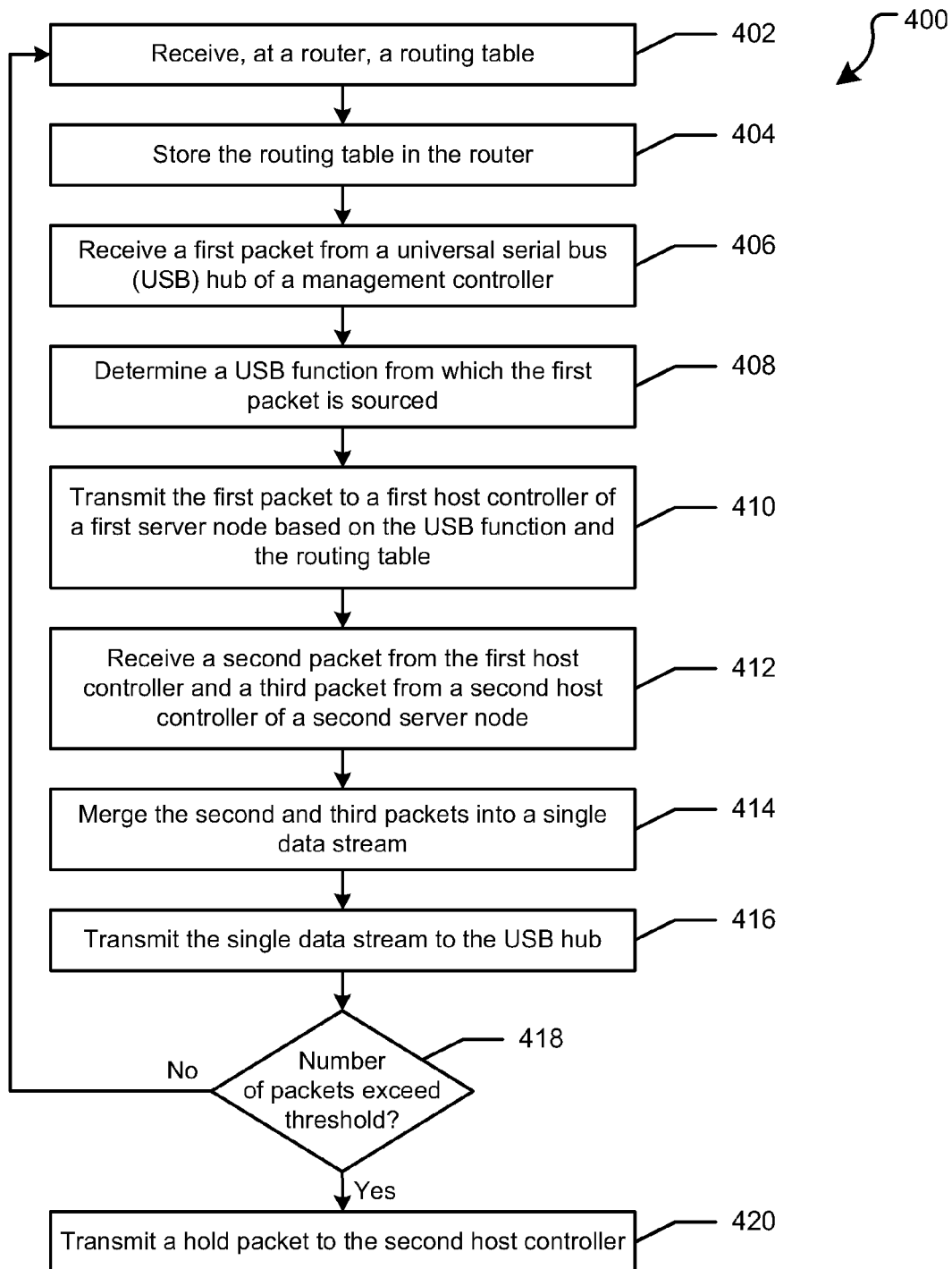


FIG. 3

**FIG. 4**

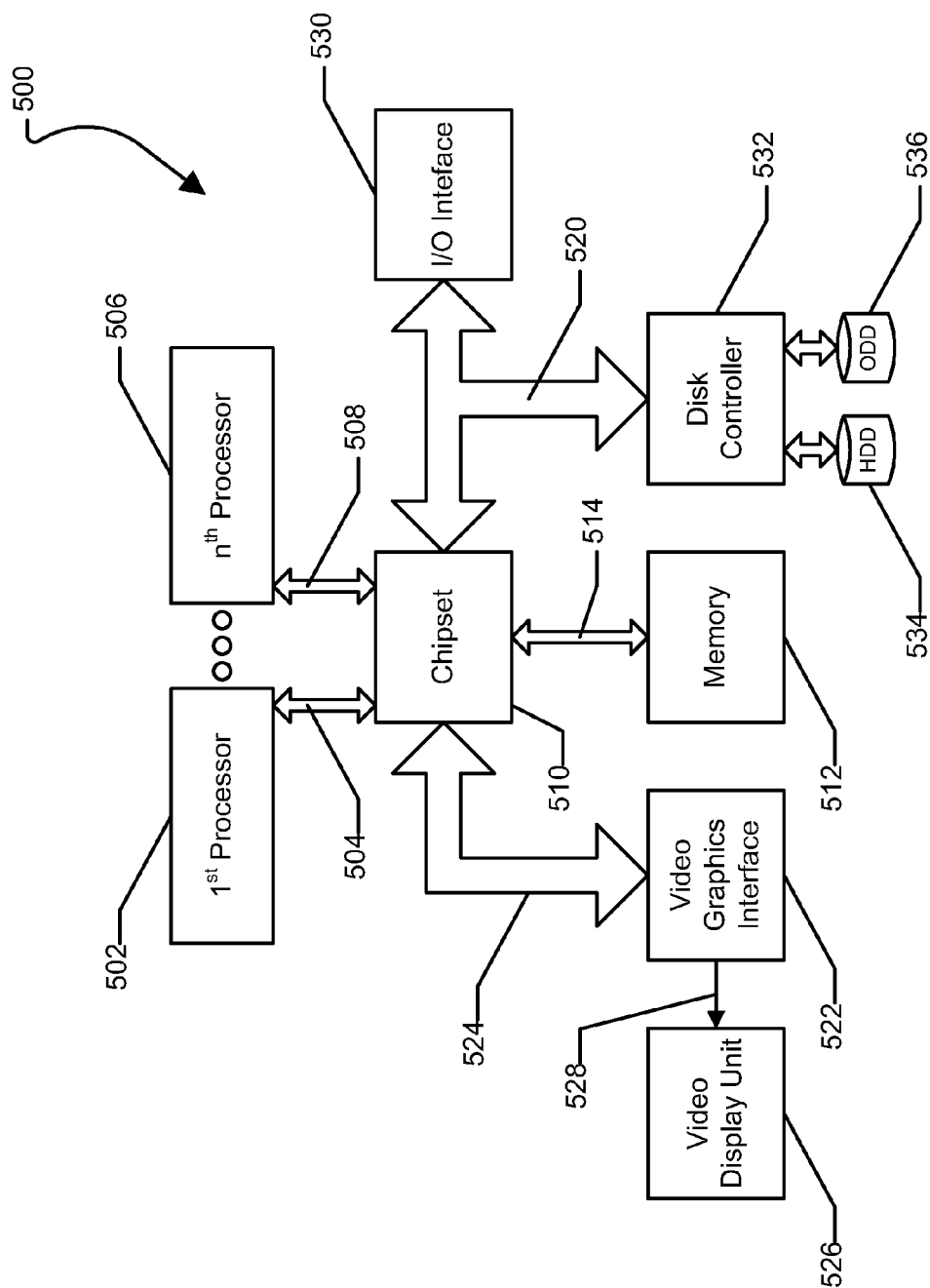


FIG. 5

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# SYSTEM AND METHOD FOR ROUTING DATA TO DEVICES WITHIN AN INFORMATION HANDLING SYSTEM

## FIELD OF THE DISCLOSURE

This disclosure generally relates to information handling systems, and more particularly relates to a system and method for routing data to devices within an information handling system.

## BACKGROUND

As the value and use of information continues to increase, individuals and businesses seek additional ways to process and store information. One option is an information handling system. An information handling system generally processes, compiles, stores, and/or communicates information or data for business, personal, or other purposes. Because technology and information handling needs and requirements can vary between different applications, information handling systems can also vary regarding what information is handled, how the information is handled, how much information is processed, stored, or communicated, and how quickly and efficiently the information can be processed, stored, or communicated. The variations in information handling systems allow for information handling systems to be general or configured for a specific user or specific use such as financial transaction processing, airline reservations, enterprise data storage, or global communications. In addition, information handling systems can include a variety of hardware and software components that can be configured to process, store, and communicate information and can include one or more computer systems, data storage systems, and networking systems. An information handling system, such as a server, can have multiple nodes or processors that can communicate with multiple devices of the information handling system.

## BRIEF DESCRIPTION OF THE DRAWINGS

It will be appreciated that for simplicity and clarity of illustration, elements illustrated in the Figures have not necessarily been drawn to scale. For example, the dimensions of some of the elements are exaggerated relative to other elements. Embodiments incorporating teachings of the present disclosure are shown and described with respect to the drawings presented herein, in which:

FIG. 1 is a block diagram of a server chassis;

FIG. 2 is an exemplary routing table for the server chassis;

FIG. 3 is a block diagram of another embodiment of the server chassis;

FIG. 4 is a flow diagram of a method for routing data packets between a universal serial bus hub and server nodes of the server chassis; and

FIG. 5 is a block diagram of a general information handling system.

The use of the same reference symbols in different drawings indicates similar or identical items.

## DETAILED DESCRIPTION OF DRAWINGS

The following description in combination with the Figures is provided to assist in understanding the teachings disclosed herein. The following discussion will focus on specific implementations and embodiments of the teachings. This focus is provided to assist in describing the teachings

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and should not be interpreted as a limitation on the scope or applicability of the teachings. However, other teachings can certainly be utilized in this application.

FIG. 1 illustrates a server chassis **100** including servers, devices, and/or information handling systems. For purposes of this disclosure, an information handling system may include any instrumentality or aggregate of instrumentalities operable to compute, classify, process, transmit, receive, retrieve, originate, switch, store, display, manifest, detect, record, reproduce, handle, or utilize any form of information, intelligence, or data for business, scientific, control, entertainment, or other purposes. For example, an information handling system may be a personal computer, a PDA, a consumer electronic device, a network server or storage device, a switch router or other network communication device, or any other suitable device and may vary in size, shape, performance, functionality, and price. The information handling system may include memory, one or more processing resources such as a central processing unit (CPU) or hardware or software control logic. Additional components of the information handling system may include one or more storage devices, one or more communications ports for communicating with external devices as well as various input and output (I/O) devices, such as a keyboard, a mouse, and a video display. The information handling system may also include one or more buses operable to transmit communications between the various hardware components.

The server chassis **100** includes a management controller **102**, a multi-host universal serial bus (USB) router **104**, and server nodes **106**, **108**, and **110**. The management controller **102** includes a USB hub **112**, an emulated control device **114**, USB physical functions **116**, **118**, **120**, and **122**, virtual USB devices **124**, **126**, **128**, and **130**, a service processor **132**, and a communication port **134**. The multi-host USB router **104** includes a packet router **136**, a packet snoopers **138**, a packet merger and arbiter **140**, a routing control module **142**, a USB host controller **144**, and host communication ports **146**, **148**, and **150**. The server nodes **106**, **108**, and **110** each include a respective USB host controller **152**, **154**, and **156**. The USB hub **112** is in communication with the USB functions **118**, **120**, and **122**, which in turn are in communication with the virtual USB devices **126**, **128**, and **130**. In an embodiment, a single USB function, such as the USB function **120**, can communicate with multiple virtual USB devices. For example, the USB function **120** can communicate with virtual USB devices **128**, **128a**, and **128b**. The USB hub **112** is also in communication with the packet router **136**, the packet snoopers **138**, and the packet merger and arbiter **140** of the multi-host USB router **104**.

The emulated control device **114** is in communication with the USB function **116**, which in turn is in communication with the virtual USB device **124**. The emulated control device **114** is also in communication with the service processor **132**, the packet router **136**, the packet snoopers **138**, and the packet merger and arbiter **140**. The packet router **136** is in communication with the USB host controller **144**. In an embodiment, the USB host controller **144** can be a virtual host controller within the multi-host USB router **104**. The packet router **136** is in communication with the USB host controllers **152**, **154**, and **156** of the server nodes **106**, **108**, and **110**. The service processor **132** is in communication with the multi-host USB router **104** via a sideband communication channel, such as an inter-integrated (I2C) channel. In different embodiments, the management controller **102** can be any type of controller for the server chassis **100**, such as an interactive Dell Remote Access controller or the like.

In an embodiment, the server chassis **100** can be a single server with multiple processors, such that each server node **106**, **108**, and **110** is a different processor within the server, and the management controller **102** can be utilized to establish different virtual connections between the server nodes and the USB devices **126**, **128**, and **130**. In another embodiment, the server chassis **100** can be a server rack and each server node **106**, **108**, and **110** can be a server that includes multiple processors. In both embodiments, each of the different server nodes **106**, **108**, and **110** only are aware of/recognize the USB hub **112** to utilize in communicating with the USB functions **118**, **120**, and **122** and the virtual USB devices **126**, **128**, and **130**. Thus, the USB packets can be streamed from the management controller **102** to the USB router **104** via the single USB bus **112** and the port **134**, and then the multi-host USB router **104** can route individual packets to each of the server nodes **106**, **108**, and **110**.

The virtual USB devices **126**, **128**, and **130** can be virtual USB keyboard video mouse (KVM) devices, virtual USB mass storage devices, virtual compact disk (CD), and/or virtual controller functions for the server nodes **106**, **108**, and **110**. The management controller **102** can dynamically assigned the different virtual USB devices **126**, **128**, and **130** to the server nodes **106**, **108**, and **110** by creating a routing table **200** as shown in FIG. 2.

FIG. 2 shows the routing table **200** including entries **202**, **204**, **206**, **208**, **210**, and **212** (202-212). In an embodiment, the assignment of the USB host controller **152**, **154**, and **156** within the routing table **200** can be based on the port that the multi-host USB router **104** is connected. The routing table **200** can indicate the USB function, the multi-host USB router port, and the type of USB device for each entry. For example, entry **202** indicates that a first function, such as USB function **116**, is assigned to a first host port **146**, and that the USB device is a virtual KVM. The entry **204** indicates that a second function, such as USB function **118**, is assigned to a second host port **148**, and that the USB device is a virtual KVM. The entry **206** indicates that a third function is assigned to the second host port, and that the USB device is a virtual CD. The entry **208** indicates that a fourth function is assigned to a third host port, and that the USB device is a lifecycle controller. The entry **210** indicates that a fifth function is assigned to a fourth host port, and that the USB device is a virtual USB mass storage device. The entry **212** indicates that a sixth function, such as USB function **120**, is assigned to the third host port, and that the USB device is a virtual USB mass storage device. The number of USB functions and USB devices within the management controller **102** can vary as devices are added or removed from the management controller. Thus, the management controller **102** can dynamically adjust the routing table **200** in response to the addition or removal of virtual USB devices, and/or based on current virtual USB devices being assigned to different USB host controllers.

Referring back to FIG. 1, the assignment of a virtual USB device **126**, **128**, or **130** to a particular server node **106**, **108**, or **110** can occur as a result of a command received by the service processor **132** from a remote user. For example, the remote user can communicate with the service processor **132** to request that a virtual KVM device be assigned to the USB host controller **152**. In this example, the service processor **132** can create the entry **202** within the routing table **200**, and the entry can assign the USB device **126** connected to the USB function **118** to the USB host controller **152** that is connected to the first host port **146** of the multi-host USB router **104**.

The remote user can also access the service processor **132** to assign other virtual USB devices in the management controller **102** as virtual devices for the USB host controllers **152**, **154**, and **156**. In an embodiment, the virtual USB devices **126**, **128**, and **130** are virtual devices of the USB host controllers **152**, **154**, and **156**, because the virtual USB devices can be dynamically assign to any one of the USB host controllers and are only a recognized device to the USB host controllers when the virtual USB device is assigned to that particular USB host controller. The service processor **132** can then send the routing table **200** to the routing control module **142** of the multi-host USB router **104**.

The assignment of a virtual USB device **126**, **128**, or **130** to a particular server node **106**, **108**, or **110** can also occur as a result of a command from the USB host controller **152**, **154**, or **156** over a low-pin count (LPC) bus between the USB host controller and the management controller **102**. For example, the USB host controller **154** can transmit a LPC bus command to the service processor **132** of the management controller **102**, and the LPC bus command can be a request to access a virtual USB mass storage device within the management controller, such as a virtual USB mass storage device on a replaceable media connected to the USB function **122**. The management controller **102** can receive the LPC bus command and then create the entry **210** in the routing table **200** assigning the virtual USB device **130** to host port four **150**, which in turn is connected to the USB host controller **156**. The service processor **132** can then send the updated routing table **200** to the routing control module **142** of the multi-host USB router **104**. In an embodiment, the service processor **132** can also receive a request for a USB device **126**, **128**, or **130** as an intelligent platform management interface (IPMI) command from one of the USB host controllers **152**, **154**, and **156**. The service processor **132** can handle the IPMI command in a substantially similar fashion as the LPC bus command described above.

The routing control module **142** can receive the routing table **200** from service processor **132** of the management controller **102**, and can store the routing table in a memory of the multi-host USB router **104**. The packet router **136** can utilize the routing control module **142** to determine a particular USB function associated with a northbound packet, and to determine a port to transmit the northbound packet. The packet router **136**, the packet snoopers **138**, and the packet merger and arbiter **140** can be utilized in the multi-host USB router **104** to transfer the USB packets to the proper USB host controller **152**, **154**, or **156**, and to send USB packets to the USB hub **112**.

For example, the multi-host USB router **104** can receive a packet from the USB function **120**. The packet snoopers **138** can snoop the packet, such as a field in the header of the packet, to determine the USB function that is the source of the packet. The routing control module **142** can then access the routing table **200** and determine the port of the multi-host USB router **104** associated with the particular USB function. For example, if the packet is from the USB function **120**, the routing table **200** can indicate that the second port is associated with that particular USB function. The packet router **136** can then transmit the packet to the USB host controller **154** via the second port **148** of the multi-host USB router **104**.

When the multi-host USB router **104** receives a USB packet from a USB host controller **152**, **154**, or **156**, the packet merger and arbiter **140** can place or merge the USB packet into a single USB data stream that can include packets from other USB host controllers, and can provide the USB data stream to the USB hub **112**, which in turn can



route the USB packet to the proper USB function 118, 120, or 122. For example, the multi-host USB router 104 can receive a USB packet from the USB host controller 152, and can receive a USB packet from the USB host controller 154. The packet merger and arbiter 140 can then merge the two USB packets into a single data stream to be provided to the USB hub 112. When the USB hub 112 receives the USB packets in the data stream, the USB hub 112 can determine that one of the packets is associated with USB function 118 and the USB device 126, and can determine that the other packet is associated with the USB function 120 and USB device 128. The USB hub 112 can then send the packets to the USB devices 126 and 128 via the respective USB functions 118 and 120.

The packet merger and arbiter 140 can utilize a time division algorithm while placing the packets from the different host controllers 152, 154, and 156 into the data stream for the USB hub 112. The time division algorithm can enable the multi-host USB router 104 to evenly distribute the packets from the different USB host controllers 152, 154, and 156 within the data stream. The multi-host USB router 104 can also determine whether a number of packets received from a particular USB host controller 152, 154, or 156 during a particular time period has exceeded a threshold number of packets. If the number of packets received from the particular USB host controller 152, 154, or 156 during the particular time period has exceeded the threshold number of packets, the multi-host USB router 104 can send a hold packet to the particular USB host controller to cause the USB host controller to stop transmitting packets for a particular amount of time.

For example, if the USB host controller 156 is performing a bulk data transfer to the virtual USB device 130, the USB host controller can send more than the threshold number of packets during a particular time period, such that the USB host controller is using all or most of the available bandwidth to the USB hub 112. The multi-host USB router 104 can then send the hold packet to the USB host controller 156 to temporarily stop the USB host controller from transmitting packets so that the other USB host controllers 152 and 154 may utilize some of the bandwidth to send packets to the virtual USB devices of the management controller 102.

In an embodiment, the packet router 136 may utilize a round-robin algorithm when transmitting packets associated with multiple KVM devices. For example, if each of the virtual USB devices 126, 128, and 130 are KVM devices that are all sending packets to the USB host controllers 152, 154, and 156 at substantially the same time, the packet router 136 can transmit one packet from each device before transmitting a second packet from one of the devices. Thus, each KVM can have even use of the bandwidth to communicate with the USB host controllers 152, 154, and 156.

In an embodiment, the routing control module 142 can include the USB host controller 144. The routing control module 142 can utilize an in-band control channel along with a USB control function, such as serial port on the management controller 102, to communicate with the management controller. In this embodiment, the routing control module 142 can scan the management controller 102 for the emulated control device 114 that is in communication with the USB function 116 and with the service processor 132. In this embodiment, the USB function 116 can be dedicated to creating the routing table, such that the virtual USB device 124 is not accessible to the server nodes 106, 108, and 110 but utilized by the service processor 132 of the management controller 102 to create the routing table. When a new USB device is added to the management controller 102, the

emulated control device 124 and the USB host controller 144 communicate determine information about the new USB device, such as whether the new USB device is a virtual KVM, a virtual USB mass storage device on a replaceable media, or the like. The USB host controller 144 and the emulated control device 114 can assign the new virtual USB device and the associated USB function to a host port of the multi-host USB router 104. The assignments can then be sent to the service processor 132 to update the routing table. The service processor 132 can then send the updated routing table to the routing control module 142 via the emulated control device 114 and the communication port 134.

FIG. 3 shows another embodiment of the server chassis 100. In this embodiment, the USB function 116 and the virtual USB device 124 can be removed from the management controller 102, and the routing control module 142 and the USB host controller 144 can be removed from the multi-host USB router 104. In this embodiment, the USB hub 112 can be removed or can be included in the management controller 102. For simplicity, the USB hub 112 has not been shown in FIG. 3. Also, the packet snoopers 128 can be replaced with an encapsulation module 302. The USB host controllers 152, 154, and 156 in the server nodes 106, 108, and 110 can communicate with the different virtual USB devices 126, 128, 128a, 128b, and 130 within the management controller 102.

In this embodiment, a USB host controller, such as the USB host controller 156, can transmit a packet 304 to a USB function and virtual USB device, such as the USB function 122 and the virtual USB device 130. The packet 304 can include a header 306 and a payload 308. The header 306 can include information to identify the USB function 122 as a destination for the packet 304, and the payload 308 can include the data for the virtual USB devices 130. When the packet 304 is received at the multi-host USB router 104 via the host port 150, the packet router 136 can determine the host port that the packet was received. The packet router can then send the packet 304 to the encapsulation module 302 along with information identifying the source port for the packet. The encapsulation module 302 can then create a new packet 310 including the header 306 from the packet 304, and a new payload 312, which can include the port identification information 314 along with the original payload 308.

The packet merger and arbiter 140 can then place or merge the new USB packet 310 into a single USB data stream that can include packets from other USB host controllers, and can provide the USB data stream to the management controller 102 via the port 134. When the packet 310 is received at the port 134, the packet can be provided to the USB function 122 based on information in the header 306 identifying the destination USB function. The service processor 132 can then send data back to the USB host controller 156.

In this situation, the service processor 132 of the management controller 102 can create the packet 310 to send data from the USB function 122 to the USB host controller 156. The packet 310 can be created with the header 306 and the payload 312, which can include a destination port identification 314 and the payload data 308. The destination port identification 314 can be generated based on the source port identification received in the management controller 102 when a corresponding packet was sent to the USB function 122. When the packet 310 is received in the multi-host USB router 104, the encapsulation module 302 can deconstruct the packet 310 and create the packet 304,

which can include the header **306** and the payload **308** that was embedded within the payload **312** of the packet **310**. The encapsulation module **302** can then provide the packet **304** to the packet router **136** with information identifying the destination port **150**. The packet router **136** can use destination port information to route the packet **304** to the USB host controller **156** via the host port **150**. Thus, in this embodiment the multi-host USB router **104** can route packets between the USB host controllers **152**, **154**, and **156** and the management controller **102** via a single USB bus connecting the multi-host USB router and the management controller without the use of a routing table.

FIG. **4** shows a method **400** for routing data packets between the USB hub **112** and the server nodes **106**, **108**, and **110** of FIG. **1**. At block **402**, a routing table is received at a router from a management controller. The routing table can identify different USB endpoint devices that are assigned or associated with different USB host controller of server nodes. The routing table is stored in the router at block **404**. In an embodiment, the routing table can be received from the management controller via a sideband channel, such as an inter-integrated circuit (I<sup>2</sup>C) channel. At block **406**, a first packet is received from a USB hub of the management controller. In an embodiment, the router is in communication with the USB hub via a single physical port of the management controller. A USB function from which the first packet is sourced is determined at block **408**. At block **410**, the first packet is transmitted to a first host controller of a first server node based on the USB function from which the packet is sourced and based on the routing table.

At block **412**, a second packet is received from the first host controller of the first server node and a third packet is received from a second host controller of a second server node. The second and third packets are merged into a single stream at block **414**. In an embodiment, the second and third packets can be merged into the single stream based on a time division algorithm. At block **416**, the single stream is transmitted to the USB hub. A determination is made whether the second host controller is transmitting more than a threshold number of packet during a specific period of time at block **418**. If the number of packets transmitted by the second host controller during the specific period of time exceeds the threshold number of packets, a hold packet is transmitted to the second host controller of the second server node at block **420**. The hold packet can ensure that the second host controller does not utilize all of the available bandwidth on the universal serial bus between the router and the USB hub. If the number of packets transmitted by the second host controller during the specific period of time does not exceed the threshold number of packets, the flow continues as stated above at block **402**.

As shown in FIG. **5**, an information handling system **500**, such as remote management station **104** or server **106**, can include a first physical processor **502** coupled to a first host bus **504** and can further include additional processors generally designated as  $n^{th}$  physical processor **506** coupled to a second host bus **508**. The first physical processor **502** can be coupled to a chipset **510** via the first host bus **504**. Further, the  $n^{th}$  physical processor **506** can be coupled to the chipset **510** via the second host bus **508**. The chipset **510** can support multiple processors and can allow for simultaneous processing of multiple processors and support the exchange of information within information handling system **500** during multiple processing operations.

According to one aspect, the chipset **510** can be referred to as a memory hub or a memory controller. For example,

the chipset **510** can include an Accelerated Hub Architecture (AHA) that uses a dedicated bus to transfer data between first physical processor **502** and the  $n^{th}$  physical processor **506**. For example, the chipset **510**, including an AHA enabled-chipset, can include a memory controller hub and an input/output (I/O) controller hub. As a memory controller hub, the chipset **510** can function to provide access to first physical processor **502** using first bus **504** and  $n^{th}$  physical processor **506** using the second host bus **508**. The chipset **510** can also provide a memory interface for accessing memory **512** using a memory bus **514**. In a particular embodiment, the buses **504**, **508**, and **514** can be individual buses or part of the same bus. The chipset **510** can also provide bus control and can handle transfers between the buses **504**, **508**, and **514**.

According to another aspect, the chipset **510** can be generally considered an application specific chipset that provides connectivity to various buses, and integrates other system functions. For example, the chipset **510** can be provided using an Intel® Hub Architecture (IHA) chipset that can also include two parts, a Graphics and AGP Memory Controller Hub (GMCH) and an I/O Controller Hub (ICH). For example, an Intel 820E, an 815E chipset, or any combination thereof, available from the Intel Corporation of Santa Clara, Calif., can provide at least a portion of the chipset **510**. The chipset **510** can also be packaged as an application specific integrated circuit (ASIC).

The information handling system **500** can also include a video graphics interface **522** that can be coupled to the chipset **510** using a third host bus **524**. In one form, the video graphics interface **522** can be an Accelerated Graphics Port (AGP) interface to display content within a video display unit **526**. Other graphics interfaces may also be used. The video graphics interface **522** can provide a video display output **528** to the video display unit **526**. The video display unit **526** can include one or more types of video displays such as a flat panel display (FPD) or other type of display device.

The information handling system **500** can also include an I/O interface **530** that can be connected via an I/O bus **520** to the chipset **510**. The I/O interface **530** and I/O bus **520** can include industry standard buses or proprietary buses and respective interfaces or controllers. For example, the I/O bus **520** can also include a Peripheral Component Interconnect (PCI) bus or a high speed PCI-Express bus. In one embodiment, a PCI bus can be operated at approximately 43 MHz and a PCI-Express bus can be operated at more than one speed, such as 2.5 GHz, 5 GHz, 8 GHz, and 16 GHz. PCI buses and PCI-Express buses can be provided to comply with industry standards for connecting and communicating between various PCI-enabled hardware devices. Other buses can also be provided in association with, or independent of, the I/O bus **520** including, but not limited to, industry standard buses or proprietary buses, such as Industry Standard Architecture (ISA), Small Computer Serial Interface (SCSI), Inter-Integrated Circuit (I<sup>2</sup>C), System Packet Interface (SPI), or Universal Serial buses (USBs).

In an alternate embodiment, the chipset **510** can be a chipset employing a Northbridge/Southbridge chipset configuration (not illustrated). For example, a Northbridge portion of the chipset **510** can communicate with the first physical processor **502** and can control interaction with the memory **512**, the I/O bus **520** that can be operable as a PCI bus, and activities for the video graphics interface **522**. The Northbridge portion can also communicate with the first physical processor **502** using first bus **504** and the second bus **508** coupled to the  $n^{th}$  physical processor **506**. The

chipset **510** can also include a Southbridge portion (not illustrated) of the chipset **510** and can handle I/O functions of the chipset **510**. The Southbridge portion can manage the basic forms of I/O such as Universal Serial Bus (USB), serial I/O, audio outputs, Integrated Drive Electronics (IDE), and ISA I/O for the information handling system **500**.

The information handling system **500** can further include a disk controller **532** coupled to the I/O bus **520**, and connecting one or more internal disk drives such as a hard disk drive (HDD) **534** and an optical disk drive (ODD) **536** such as a Read/Write Compact Disk (R/W CD), a Read/Write Digital Video Disk (R/W DVD), a Read/Write mini-Digital Video Disk (R/W mini-DVD), or other type of optical disk drive.

Although only a few exemplary embodiments have been described in detail in the exemplary embodiments without materially departing from the novel teachings and advantages of the embodiments of the present disclosure. For example, the methods described in the present disclosure can be stored as instructions in a computer readable medium to cause a processor, such as chipset **510**, to perform the method. Additionally, the methods described in the present disclosure can be stored as instructions in a non-transitory computer readable medium, such as a hard disk drive, a solid state drive, a flash memory, and the like. Accordingly, all such modifications are intended to be included within the scope of the embodiments of the present disclosure as defined in the following claims. In the claims, means-plus-function clauses are intended to cover the structures described herein as performing the recited function and not only structural equivalents, but also equivalent structures.

What is claimed is:

1. An information handling system comprising:
  - a management controller including a universal serial bus hub and first and second devices, the management controller to assign the first device to a first server node, to assign the second device to a second server node, and to create a routing table associated with the assignment of the first and second devices respectively to the first and second server nodes, wherein the first and second server nodes are only aware of the universal serial bus hub for communication with the management controller; and
  - a router in communication with the management controller, the router to receive the routing table from the management controller, to receive data from the first server node, and to route the data to the first device based on the routing table, wherein the universal serial bus hub communicates with the router via a single physical port of the management controller.
2. The information handling system of claim 1 wherein the management controller further includes a service processor, the service processor to receive a request to assign the first device to the first server node.
3. The information handling system of claim 2 wherein the request is received in the service processor via a remote user initiated transaction.
4. The information handling system of claim 1 wherein the management controller further includes a service processor, the service processor to determine an assignment of the first device to the first server node, and to assign the first device to the first server node in the routing table.
5. The information handling system of claim 2 wherein the request is received in the service processor via a command from the first server node.
6. The information handling system of claim 1 wherein the router merges data from the first and second server nodes

into a single stream that is transmitted to the universal serial bus hub via the single physical port.

7. The information handling system of claim 1 wherein the first and second device are virtual universal serial bus devices within the management controller.

8. An information handling system comprising:

first and second server nodes each including a host controller;

a management controller including first and second devices, the management controller to assign the first device to the first server node, to assign the second device to the second server node, and to create a first packet to send data from the first device to the first server node, the first packet including a header including header fields, and a first payload having a destination port identification and a first embedded payload; and

a router in communication with the server nodes and with the management controller, the router to receive the first packet from the management controller, to deconstruct the first payload of the first packet, and to create a second packet including the header and the first embedded payload of the first packet, and to send the second packet to the first server node via a first port of the router based on the destination port identification in the first packet.

9. The information handling system of claim 8 wherein the router is further configured to receive a third packet from the second server node, the third packet including a second payload having data for the second device, to determine a second port of the router as a source port for the third packet, to create a fourth packet, the fourth packet including a third payload having the second payload as a second embedded payload and a source identification to identify the second port, and to send the fourth packet to the management controller.

10. The information handling system of claim 9 wherein the router merges the fourth packet with packets from the first server node into a single stream that is sent to the management controller via a single communication bus between the router and the management controller.

11. The information handling system of claim 8 wherein the first and second device are virtual universal serial bus devices within the management controller.

12. The information handling system of claim 8 wherein the first and second server nodes are processors within the information handling system.

13. The information handling system of claim 8 wherein the first and second server nodes are separate servers within the information handling system, and wherein each of the first and second servers includes one or more processors.

14. A method comprising:

receiving, at a router, a routing table from a management controller, wherein the routing table is received via a sideband communication;

receiving, at the router, a first packet from a universal serial bus hub of the management controller;

determining a universal serial bus function from which the first packet is sourced;

transmitting the first packet to a first host controller of a first server node based on the universal serial bus function from which the first packet is sourced and based on the routing table;

receiving a second packet from the first host controller of the first server node and a third packet from a second host controller of a second server node;

merging the second and third packets into a single stream;  
and  
transmitting the single stream to the universal serial bus  
hub.

**15.** The method of claim **14** further comprising: 5  
transmitting a hold packet to the second host controller of  
the second server node in response to the second host  
controller transmitting a number of packets that  
exceeds a threshold number packets during a specific  
period of time. 10

**16.** The method of claim **14** wherein the second and third  
packets are merged into the single stream based on a time  
division algorithm.

**17.** The method of claim **14** wherein the router is in  
communication with the universal serial bus hub via a single 15  
physical port of the management controller.

**18.** The method of claim **14** wherein the routing table is  
received from the management controller via a sideband  
channel.

**19.** The method of claim **14** wherein the first and second 20  
server nodes are processors within the information handling  
system.

**20.** The method of claim **14** wherein the first and second  
server nodes are separate servers within the information  
handling system, and wherein each of the first and second 25  
servers includes one or more processors.

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